

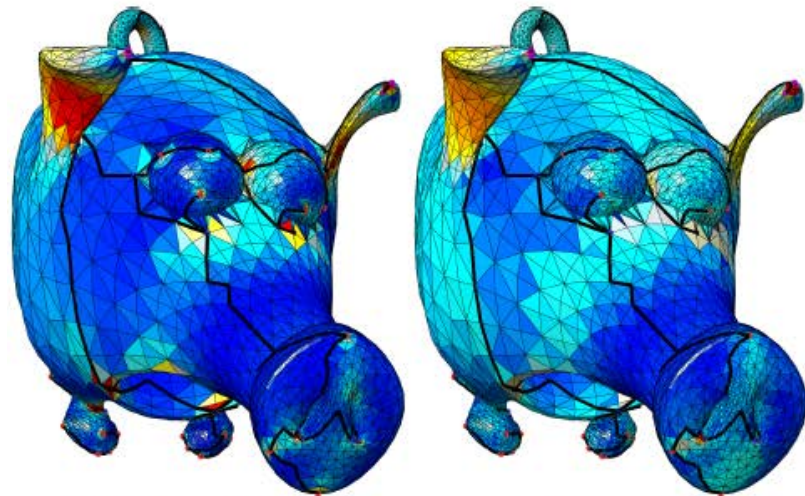
# Strict Minimizers For Geometric Optimization

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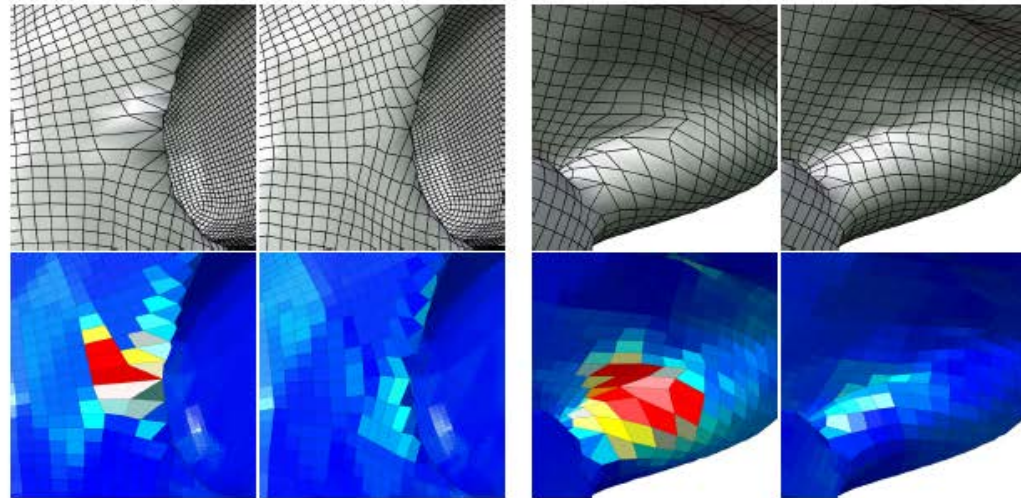
# Motivation

[Aigerman and Lipman 2013] vs. our strict minimizer



*bounded distortion*

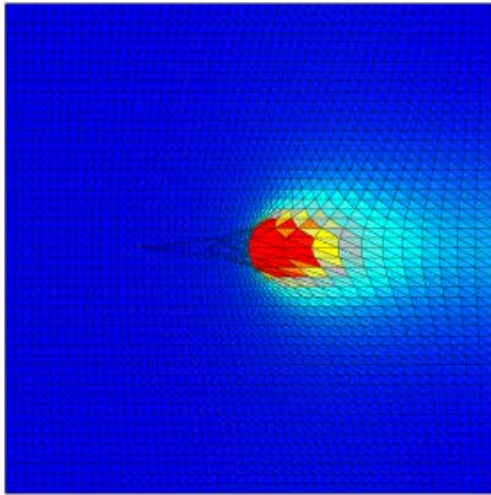
*strict minimizer*



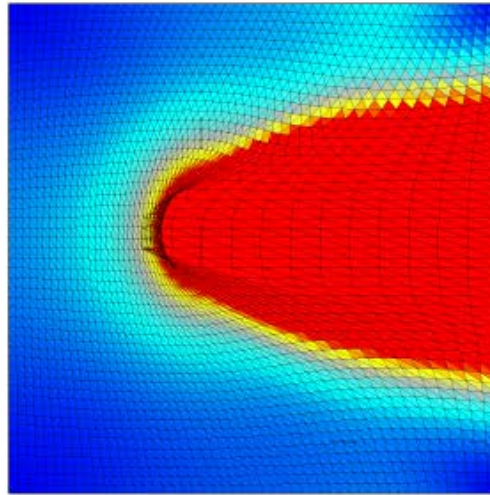
*bounded distortion* *strict minimizer*

*bounded distortion* *strict minimizer*

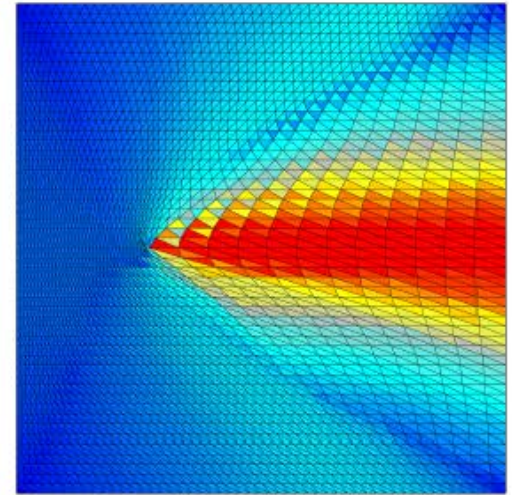
# A toy example



$L_2$



$L_\infty$



*strict*

# Overall idea

1. Find  $L_\infty$  solution with smallest max distortion set of facets
2. Fix the distortion on facets with max distortion
3. Repeat until all facets are fixed

# Questions

- Can we define “the best”  $L_\infty$  solution independent of the algorithm?
- Is it unique?
- Is there an efficient algorithm?

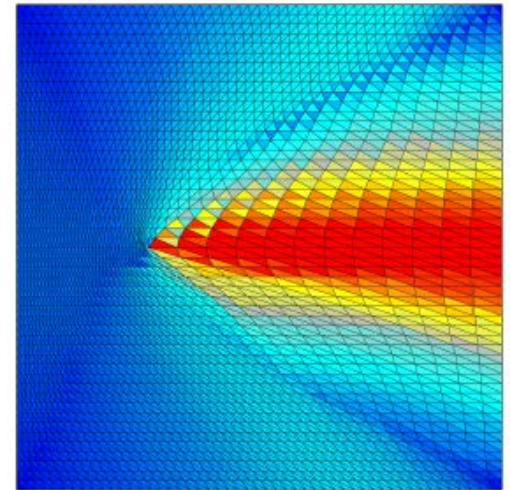
# Questions

- Can we define “the best”  $L_\infty$  solution independent of the algorithm?
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# Lexicographical ordering

Vector of facets' distortion for a mapping  $f$

$$D[\mathbf{f}] = [D_1[\mathbf{f}], \dots, D_N[\mathbf{f}]]$$



$D_i(f)$  is the distortion on facet  $i$ .

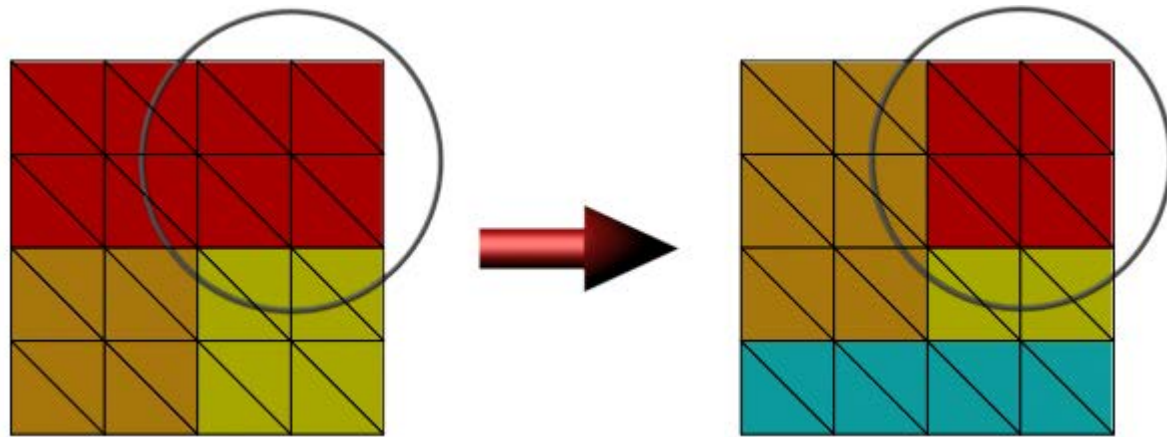
Define an ordering based on this (sorted) vector

(For simplicity, we consider a mesh with similar facets.)

(Lexicographical is a dictionary order: “3322” < “3331”)

# How to compute this in theory

1. Solve  $L_\infty$
2. Figure out minimal subset of facets that must have max distortion
3. Fix the distortion on facets with max distortion
4. Repeat until all facets are fixed

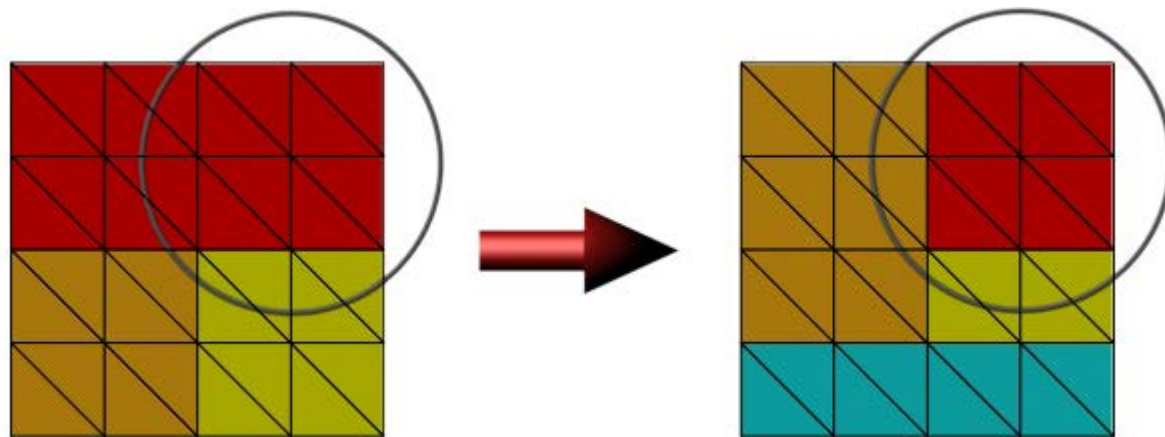


# Minimize the distortion of free facets

Consider the set of solutions  $X$ :

$$k = \min \max_{j \notin I} D_j[\mathbf{f}], \text{ subject to } D_i[\mathbf{f}] \leq k_i, i \in I$$

$I$  is the set of facets already constrained to have an upper bound  $k_i$  on the distortion.





# Questions

- Can we define “the best”  $L_\infty$  solution independent of the algorithm?
- Is it unique?
- Is there an efficient algorithm?

# Essential facets

Facet  $j$  is essential if for all  $L_\infty$  solutions in  $X$ .

$$D_j[\mathbf{f}] = k$$

**Proposition 1.** The essential set is not empty.

The essential set is unique by definition.

# Theoretical algorithm

1. Solve  $L_\infty$
2. Find the essential set
3. Add essential set to the fixed set
4. Repeat until all facets are fixed

# Uniqueness

**Proposition 2.** Any strict minimizer has the same distortion vector

$$D[\mathbf{f}^*] = [k_1, \dots, k_N]$$

Uniqueness of solutions vs. vectors of distortion

# Questions

- Can we define “the best”  $L_\infty$  solution independent of the algorithm?
- Is it unique?
- Is there an efficient algorithm?

# $L_\infty$ optimization

min  $k$ , subject to

$$D_i^{FD}(\mathbf{q}) = \|J_i(\mathbf{q}) - R_i\|_F \leq k, \quad i = 1 \dots N$$

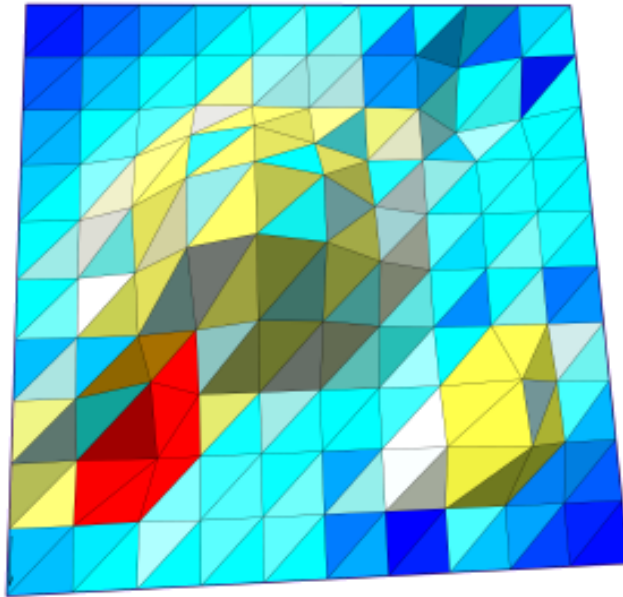
- Isometric or quasiconformal (ASAP)
- A SOCP problem, which can be solve efficiently
- Always feasible
- Can operate inside a local/global framework
- Distortion below 1 means no foldovers

# Theoretical algorithm

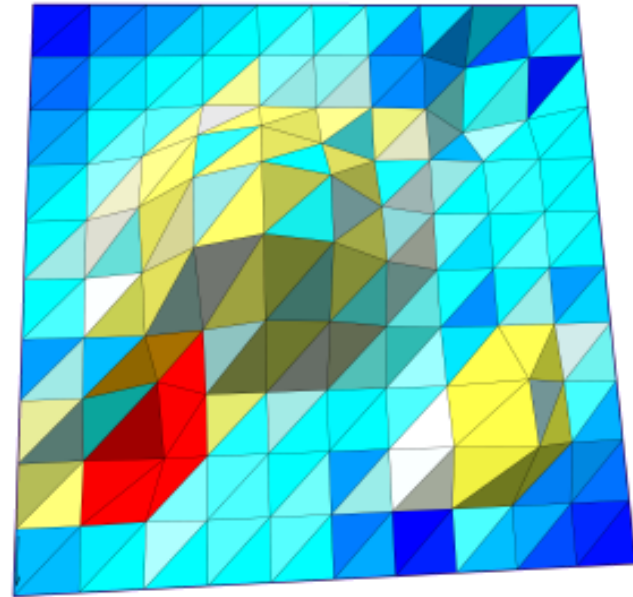
1. Solve  $L_\infty$
2. Find the essential set
3. Add essential set to the fixed set
4. Repeat until all facets are fixed

# First approximate algorithm

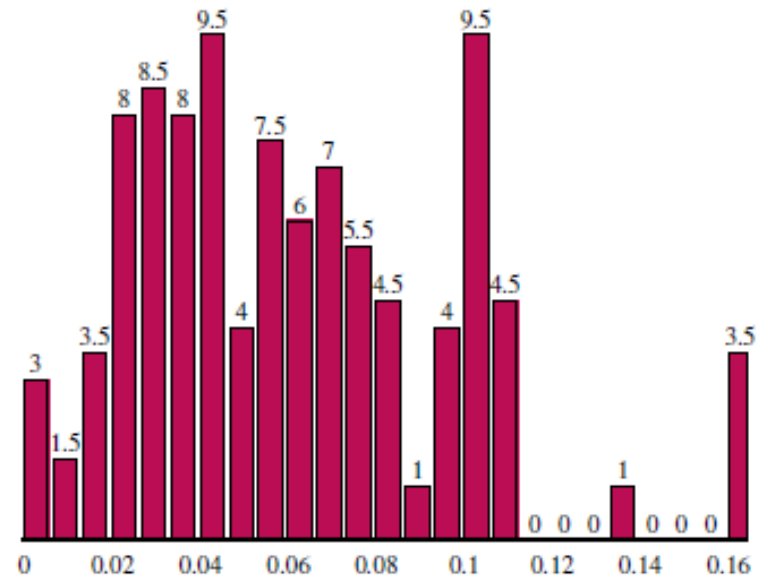
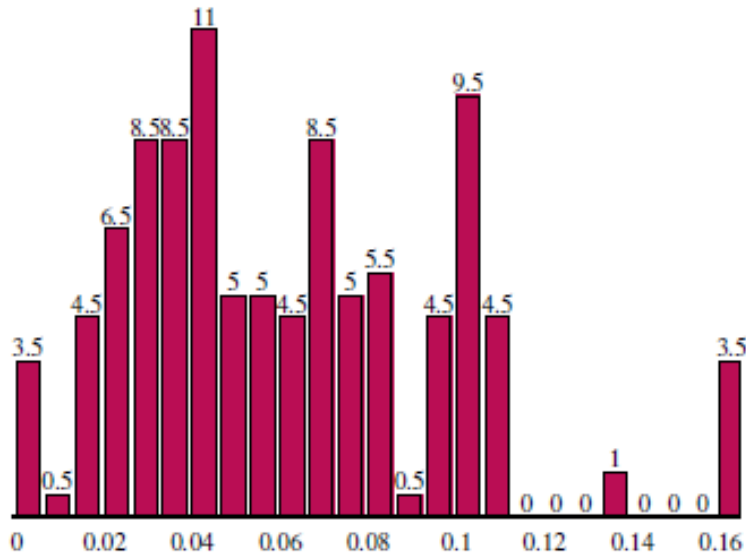
1. Solve  $L_\infty$  for upper bound  $k$  on free facets
2. Solve  $L_2$ , temporarily constraining free facets to  $k$
3. Estimate essential set as facets with distortion  $k-\Delta$
4. Add essential set to the fixed set
5. Repeat until all facets are fixed



Approximate



Exact





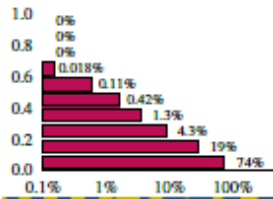
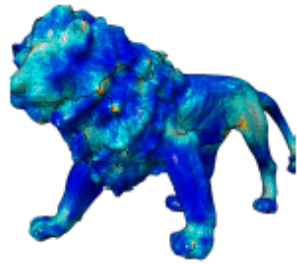
# Second approximate algorithm

- Solve  $L_\infty$  to obtain an initial solution.
- **repeat** 50 times
  - **for** each vertex  $v$ 
    - Solve 1-ring  $L_\infty$  problem on  $v$

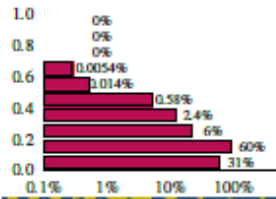
## Properties

- Coordinate descent
- One vertex at a time
- Problem can be solved semi-explicitly

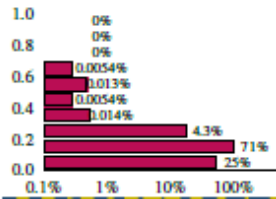
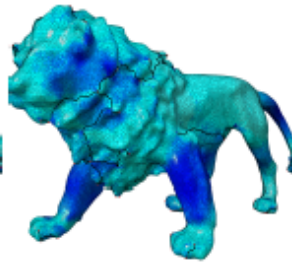
$L_\infty$ -minimizer



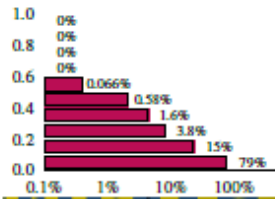
strict, relax



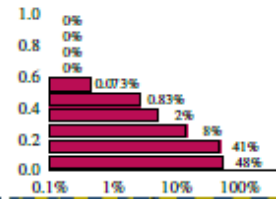
strict, direct



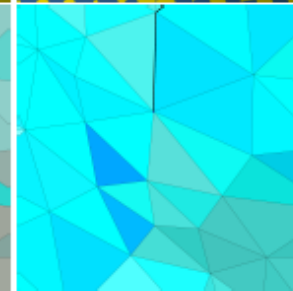
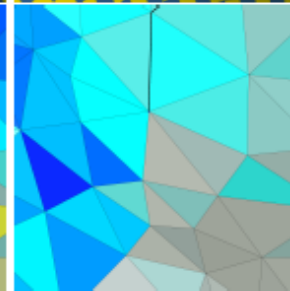
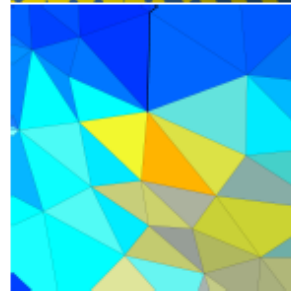
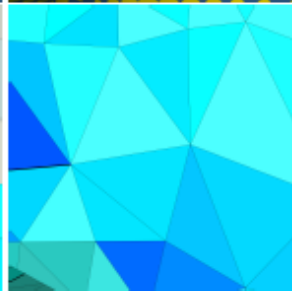
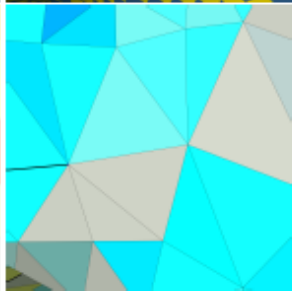
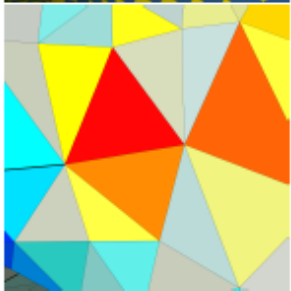
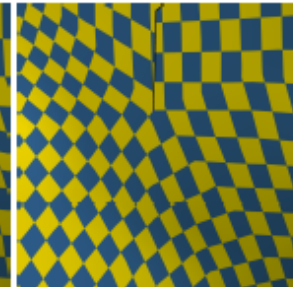
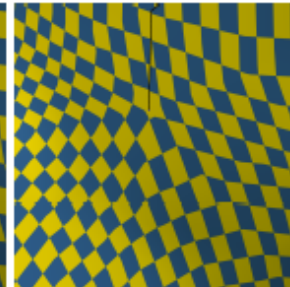
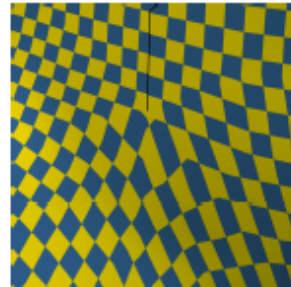
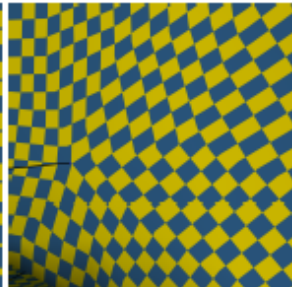
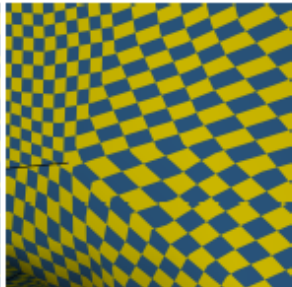
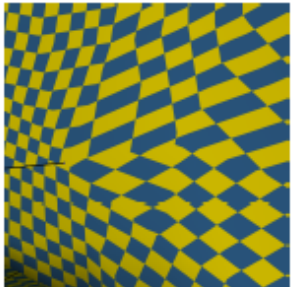
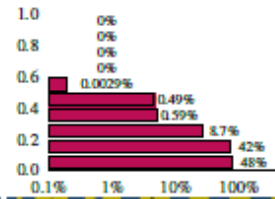
$L_\infty$ -minimizer



strict, relax



strict, direct



[Bommes et al. 2013]

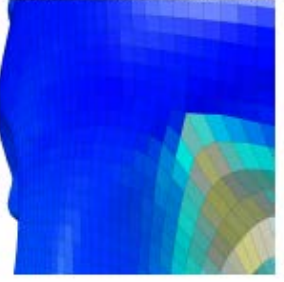
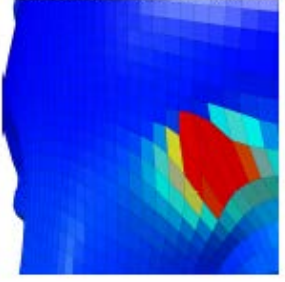
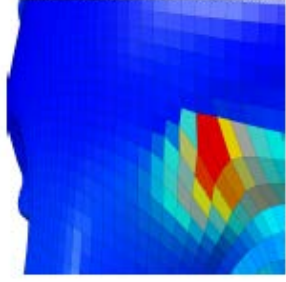
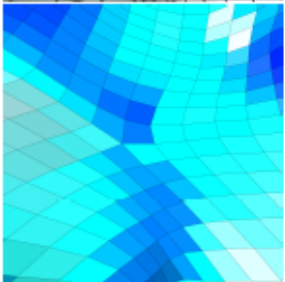
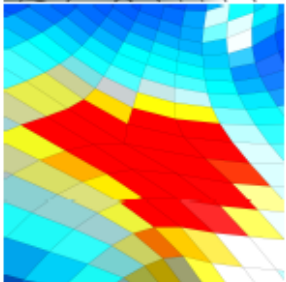
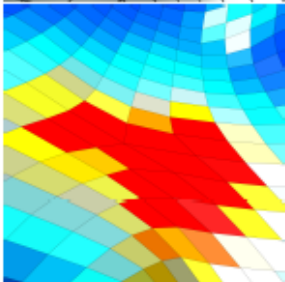
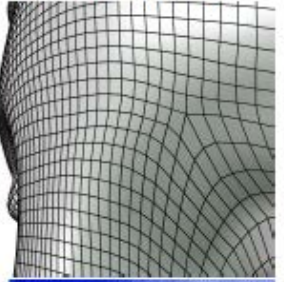
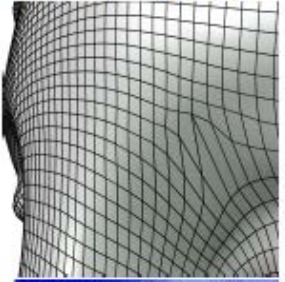
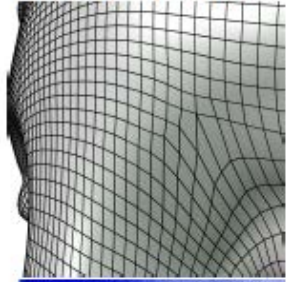
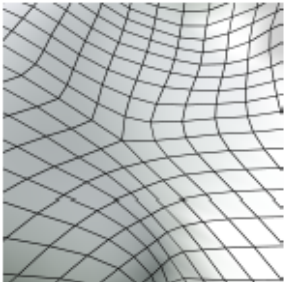
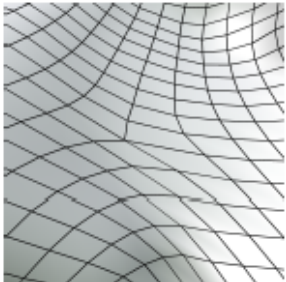
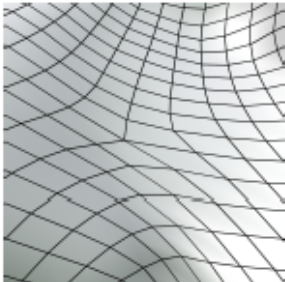
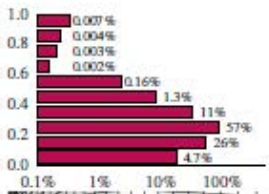
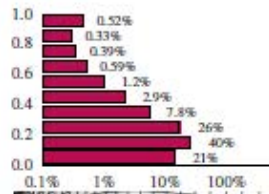
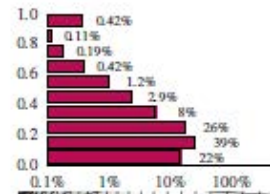
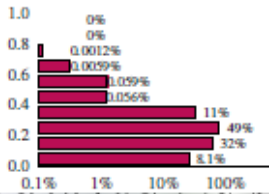
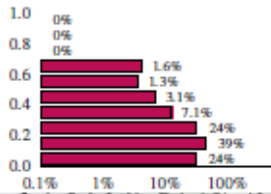
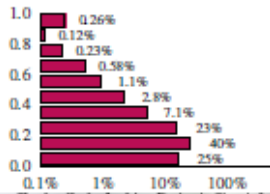
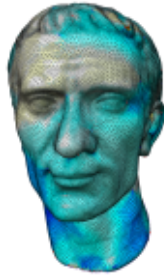
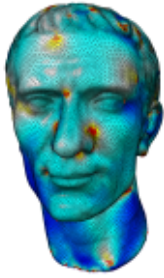
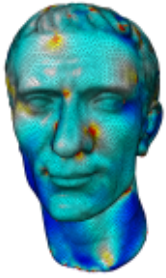
$L_\infty$ -minimizer

strict minimizer

[Bommes et al. 2013]

$L_\infty$ -minimizer

strict minimizer



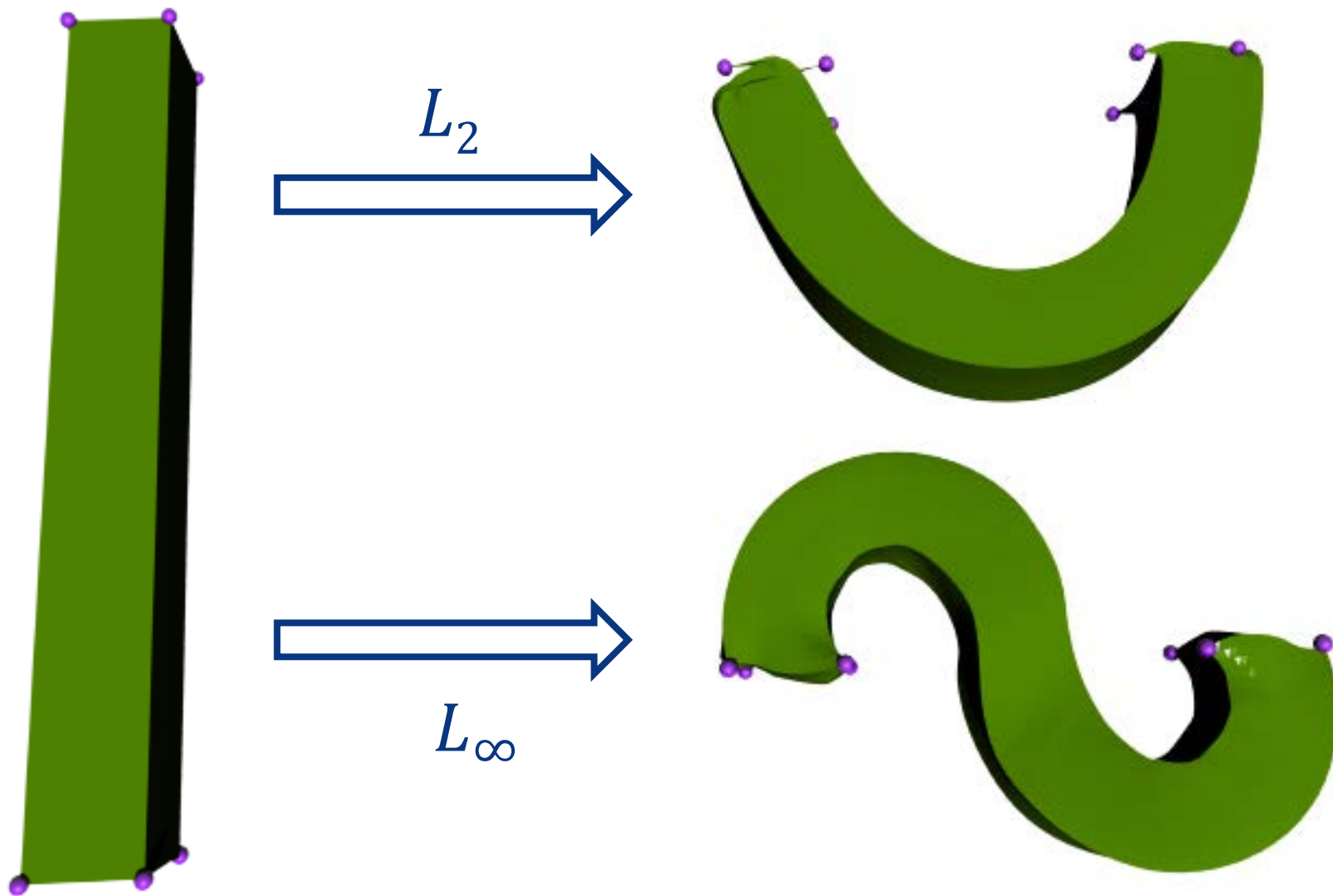


$L_\infty$

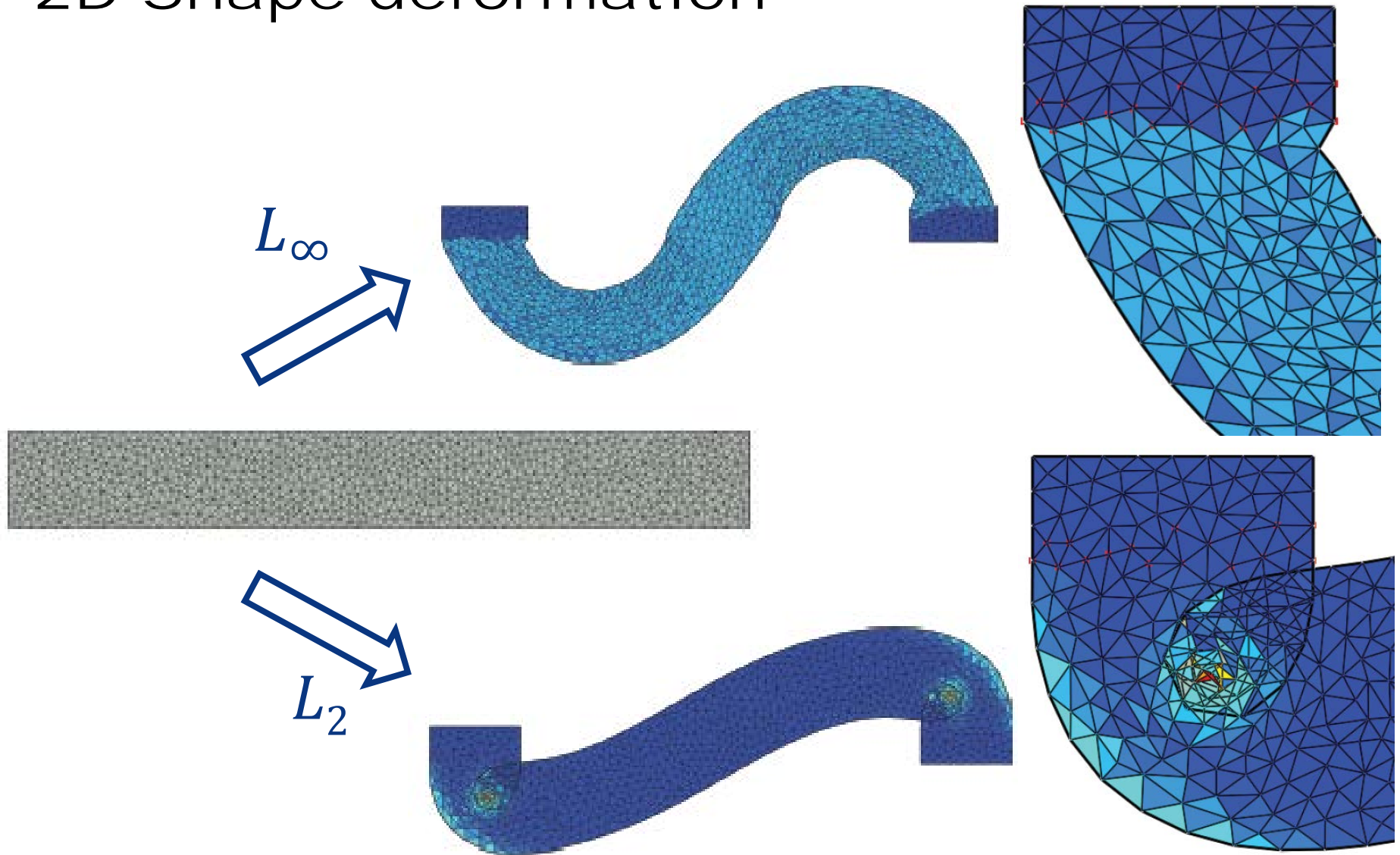


strict min

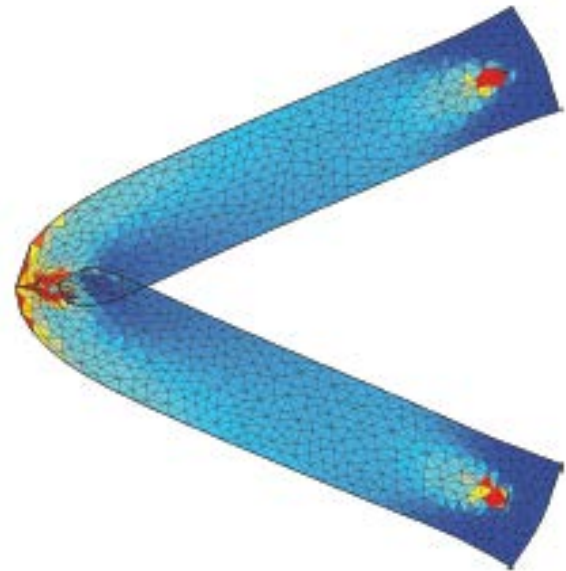
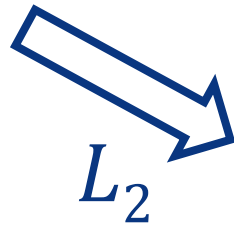
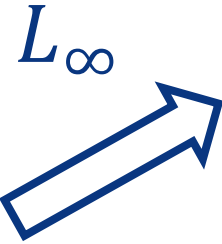
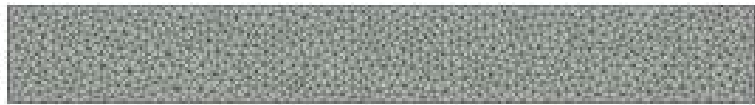
# 3D Shape deformation



# 2D Shape deformation



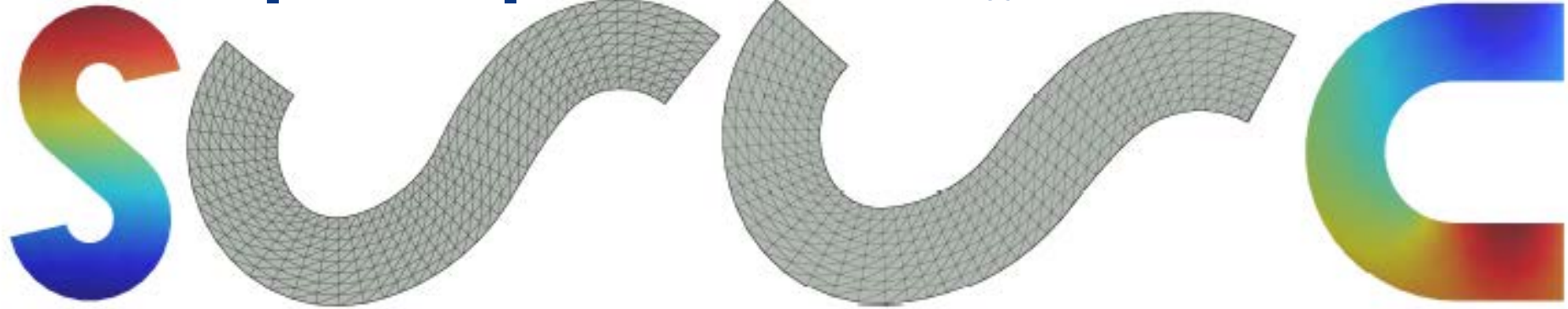
# 2D Shape deformation



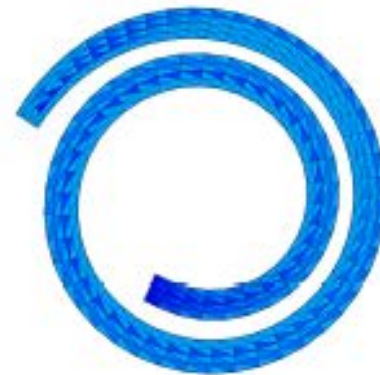
# Shape interpolation

[Chen13]

$L_\infty$



$L_2$

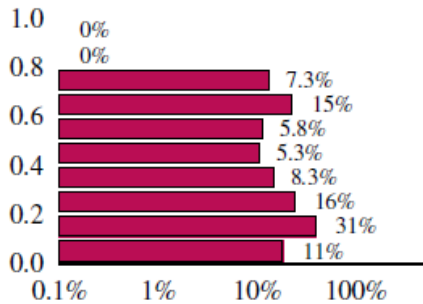
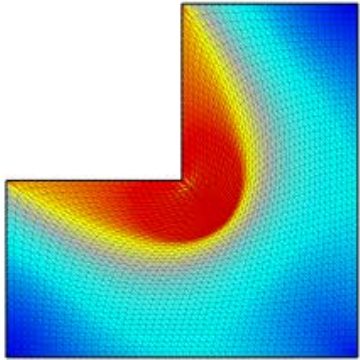
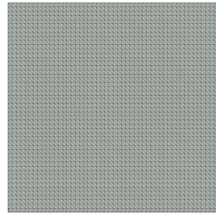


$L_\infty$

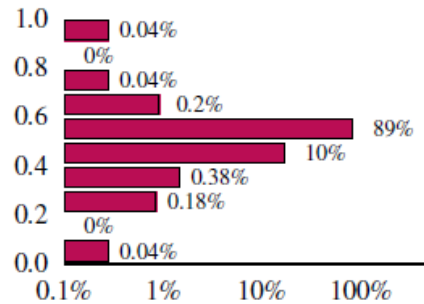
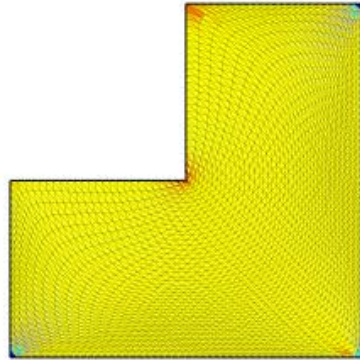


# Teichmuller

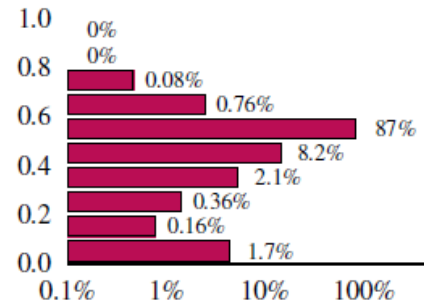
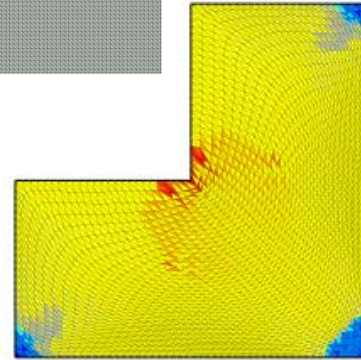
Source



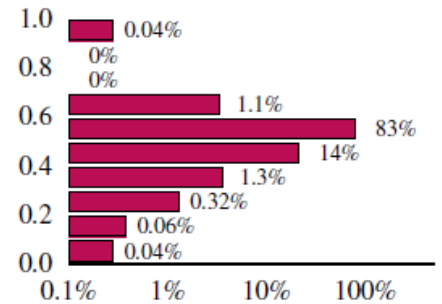
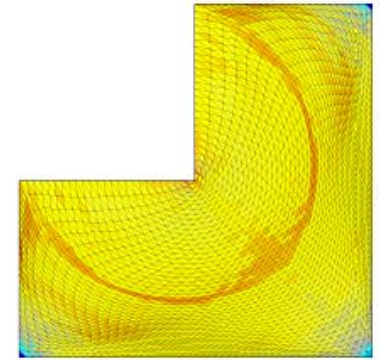
[Lipman12]



[Weber12]



Strict min [Lipman12]



Strict min ASAP

## • Conclusions

- Defined strict minimizer – ‘best’  $L_\infty$  solution
- Uniqueness
- Efficient approximate algorithms

## • Acknowledgments

- Julian Panetta
- NSF awards IIS-1320635 and IIS-1247240